

60  
JCS74 U.S. PTO  
69/24/99

Please type a plus sign (+) inside this box [+]

PTO/SB/05 (12/97)

Approved for use through 09/30/00. OMB 0651-0032

Patent and Trademark Office: U.S. DEPARTMENT OF COMMERCE  
Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number.

JCS50 U.S. PTO  
69/24/99

## UTILITY PATENT APPLICATION TRANSMITTAL

(Only for new nonprovisional applications under 37 CFR 1.53(b))

Attorney Docket No. 003498.P033

Total Pages 5

First Named Inventor or Application Identifier Rajugopal R. Gubbi

Express Mail Label No. EL371011441US

**ADDRESS TO:** Assistant Commissioner for Patents  
Box Patent Application  
Washington, D. C. 20231

### APPLICATION ELEMENTS

See MPEP chapter 600 concerning utility patent application contents.

1.  Fee Transmittal Form  
(Submit an original, and a duplicate for fee processing)
2.  Specification (Total Pages 24)  
(preferred arrangement set forth below)
  - Descriptive Title of the Invention
  - Cross References to Related Applications
  - Statement Regarding Fed sponsored R & D
  - Reference to Microfiche Appendix
  - Background of the Invention
  - Brief Summary of the Invention
  - Brief Description of the Drawings (if filed)
  - Detailed Description
  - Claims
  - Abstract of the Disclosure
3.  Drawings(s) (35 USC 113) (Total Sheets 2)
4.  Oath or Declaration (Total Pages 4)
  - a.  Newly Executed (Original or Copy)
  - b.  Copy from a Prior Application (37 CFR 1.63(d))  
(for Continuation/Divisional with Box 17 completed) (Note Box 5 below)
  - i.  DELETIONS OF INVENTOR(S) Signed statement attached deleting inventor(s) named in the prior application, see 37 CFR 1.63(d)(2) and 1.33(b).
5.  Incorporation By Reference (useable if Box 4b is checked)  
The entire disclosure of the prior application, from which a copy of the oath or declaration is supplied under Box 4b, is considered as being part of the disclosure of the accompanying application and is hereby incorporated by reference therein.
6.  Microfiche Computer Program (Appendix)

7.  Nucleotide and/or Amino Acid Sequence Submission  
(if applicable, all necessary)  
a.  Computer Readable Copy  
b.  Paper Copy (identical to computer copy)  
c.  Statement verifying identity of above copies

#### ACCOMPANYING APPLICATION PARTS

8.  Assignment Papers (cover sheet & documents(s))  
9.  a. 37 CFR 3.73(b) Statement (where there is an assignee)  
 b. Power of Attorney  
10.  English Translation Document (if applicable)  
11.  a. Information Disclosure Statement (IDS)/PTO-1449  
 b. Copies of IDS Citations  
12.  Preliminary Amendment  
13.  Return Receipt Postcard (MPEP 503) (Should be specifically itemized)  
14.  a. Small Entity Statement(s)  
 b. Statement filed in prior application, Status still proper and desired  
15.  Certified Copy of Priority Document(s) (if foreign priority is claimed)  
16.  Other: Copy of Postcard w/Express Mail Stamp

17. If a **CONTINUING APPLICATION**, check appropriate box and supply the requisite information:

Continuation  Divisional  Continuation-in-part (CIP)  
of prior application No: \_\_\_\_\_

18. **Correspondence Address**

Customer Number or Bar Code Label

(Insert Customer No. or Attach Bar Code Label here)

or

Correspondence Address Below

NAME Tarek N. Fahmi – Reg. No.: 41,402

BLAKELY, SOKOLOFF, TAYLOR & ZAFMAN LLP



ADDRESS 12400 Wilshire Boulevard

Seventh Floor

CITY Los Angeles STATE California ZIP CODE 90025-1026

Country U.S.A. TELEPHONE (408) 720-8598 FAX (408) 720-9397

12/01/97

- 2 -

PTO/SB/05 (12/97)

Approved for use through 09/30/00. OMB 0651-0032  
Patent and Trademark Office; U.S. DEPARTMENT OF COMMERCE

UNITED STATES PATENT APPLICATION

For

**METHOD AND APPARATUS FOR ACCOMMODATING ASYNCHRONOUS  
DATA TRANSMISSIONS IN A WIRELESS COMPUTER NETWORK**

Inventors:

Rajugopal R. Gubbi

Prepared by:

BLAKELY SOKOLOFF TAYLOR & ZAFMAN LLP  
12400 Wilshire Boulevard  
Los Angeles, CA 90025-1026  
(408) 720-8598

Attorney's Docket No.: 003498.P033

"Express Mail" mailing label number: EL371011441US

Date of Deposit: September 24, 1999

I hereby certify that I am causing this paper or fee to be deposited with the United States Postal Service "Express Mail Post Office to Addressee" service on the date indicated above and that this paper or fee has been addressed to the Assistant Commissioner for Patents, Washington, D. C. 20231

Patricia A. Balero

(Typed or printed name of person mailing paper or fee)

Patricia A. Balero  
(Signature of person mailing paper or fee)

9/24/99

(Date signed)

## **METHOD AND APPARATUS FOR ACCOMMODATING ASYNCHRONOUS DATA TRANSMISSIONS IN A WIRELESS COMPUTER NETWORK**

### **RELATED APPLICATION**

5        This application is a continuation-in-part of co-pending Application No. 09/151,579, entitled "Method and Apparatus for Accessing a Computer Network Communication Channel", filed September 11, 1998, by Rajugopal R. Gubbi, Natarajan Ekambaram and Nirmalendu Bikash Patra, and assigned to the Assignee of the present application.

10

### **FIELD OF THE INVENTION**

The present invention relates generally to a scheme for communications within a computer network and, in particular, to asynchronous communications as occur between a central server and one or more clients across a wireless communications link.

15

### **BACKGROUND**

Modern computer networks allow for inter-communication between a number of nodes such as personal computers, workstations, peripheral units and the like. Network links transport information between these nodes, which may sometimes be separated by large distances. However, to date most computer networks have relied on wired links to transport this information. Where wireless links are used, they have typically been components of a very large network, such as a wide area network, which may employ satellite communication links to interconnect network nodes separated by very large distances. In such cases, the transmission protocols used across the wireless links have generally been established by the service entities carrying the data being transmitted, for example, telephone companies and other service providers.

In the home environment, computers have traditionally been used as stand-alone devices. More recently, however, there have been some steps taken to integrate the home computer with other appliances. For example, in so-called "Smart Homes", computers may be used to turn on and off various appliances and to control their operational settings.

5 In such systems, wired communication links are used to interconnect the computer to the appliances that it will control. Such wired links are expensive to install, especially where they are added after the original construction of the home.

In an effort to reduce the difficulties and costs associated with wired communication links, some systems for interconnecting computers with appliances have

10 utilized analog wireless links for transporting information between these units. Such analog wireless links operate at frequencies commonly utilized by wireless telephones. Although easier to install than conventional wired communication links, analog wireless communication links suffer from a number of disadvantages. For example, degraded signals may be expected on such links because of multipath interference. Furthermore,

15 interference from existing appliances, such as televisions, cellular telephones, wireless telephones and the like may be experienced. Thus, analog wireless communication links offer less than optimum performance for a home environment.

In the above-referenced co-pending application, Serial No. 09/151,579, which is incorporated herein by reference, a computer network employing a digital, wireless

20 communication link adapted for use in the home environment was described. That architecture included a number of network components arranged in a hierarchical fashion and communicatively coupled to one another through communication links operative at different levels of the hierarchy. At the highest level of the hierarchy, a communication protocol that supports dynamic addition of new network components at any level of the

hierarchy according to bandwidth requirements within a communication channel operative at the highest level of the network hierarchy is used.

The generalization of this network structure is shown in **Figure 1**. A subnet 10 includes a server 12. In this scheme, the term "subnet" is used to describe a cluster of 5 network components that includes a server and several clients associated therewith (e.g., coupled through the wireless communication link). Depending on the context of the discussion however, a subnet may also refer to a network that includes a client and one or more subclients associated therewith. A "client" is a network node linked to the server through the wireless communication link. Examples of clients include audio/video 10 equipment such as televisions, stereo components, personal computers, satellite television receivers, cable television distribution nodes, and other household appliances.

Server 12 may be a separate computer that controls the communication link, however, in other cases server 12 may be embodied as an add-on card or other component attached to a host computer (e.g., a personal computer) 13. Server 12 has an associated 15 radio 14, which is used to couple server 12 wirelessly to the other nodes of subnet 10. The wireless link generally supports both high and low bandwidth data channels and a command channel. Here a channel is defined as the combination of a transmission frequency (more properly a transmission frequency band) and a pseudo-random (PN) code used in a spread spectrum communication scheme. In general, a number of available 20 frequencies and PN codes may provide a number of available channels within subnet 10. As is described in the co-pending application cited above, servers and clients are capable of searching through the available channels to find a desirable channel over which to communicate with one another.

Also included in subnet 10 are a number of clients 16, some of which have shadow 25 clients 18 associated therewith. A shadow client 18 is defined as a client which receives

the same data input as its associated client 16 (either from server 12 or another client 16), but which exchanges commands with server 12 independently of its associated client 16.

Each client 16 has an associated radio 14, which is used to communicate with server 12, and some clients 16 may have associated subclients 20. Subclients 20 may include

5 keyboards, joysticks, remote control devices, multi-dimensional input devices, cursor control devices, display units and/or other input and/or output devices associated with a particular client 16. A client 16 and its associated subclients 20 may communicate with one another via communication links 21, which may be wireless (e.g., infra-red, ultrasonic, spread spectrum, etc.) communication links.

10 Each subnet 10 is arranged in a hierarchical fashion with various levels of the hierarchy corresponding to levels at which intra-network component communication occurs. At a highest level of the hierarchy exists the server 12 (and/or its associated host 13), which communicates with various clients 16 via the wireless radio channel. At other, lower levels of the hierarchy the clients 16 communicate with their various subclients 20  
15 using, for example, wired communication links or wireless communication links such as infrared links.

Where half-duplex radio communication is used on the wireless link between server 12 and clients 16, a communication protocol based on a slotted link structure with dynamic slot assignment is employed. Such a structure supports point-to-point connections within 20 subnet 10 and slot sizes may be re-negotiated within a session. Thus a data link layer that supports the wireless communication can accommodate data packet handling, time management for packet transmission and slot synchronization, error correction coding (ECC), channel parameter measurement and channel switching. A higher level transport layer provides all necessary connection related services, policing for bandwidth utilization, 25 low bandwidth data handling, data broadcast and, optionally, data encryption. The

transport layer also allocates bandwidth to each client 16, continuously polices any under or over utilization of that bandwidth, and also accommodates any bandwidth renegotiations, as may be required whenever a new client 16 comes on-line or when one of the clients 16 (or an associated subclient 20) requires greater bandwidth.

5        The slotted link structure of the wireless communication protocol for the transmission of real time, multimedia data (e.g., as frames) within a subnet 10 is shown in **Figure 2**. At the highest level within a channel, forward (F) and backward or reverse (B) slots of fixed (but negotiable) time duration are provided within each frame transmission period. During forward time slots F, server 12 may transmit video and/or audio data

10      and/or commands to clients 16, which are placed in a listening mode. During reverse time slots B, server 12 listens to transmissions from the clients 16. Such transmissions may include audio, video or other data and/or commands from a client 16 or an associated subclient 20. At the second level of the hierarchy, each transmission slot (forward or reverse) is made up of one or more radio data frames 40 of variable length. Finally, at the

15      lowest level of the hierarchy, each radio data frame 40 is comprised of server/client data packets 42, which may be of variable length.

Each radio data frame 40 is made up of one server/client data packet 42 and its associated error correction coding (ECC) bits. Variable length framing is preferred over constant length framing in order to allow smaller frame lengths during severe channel

20      conditions and vice-versa. This adds to channel robustness and bandwidth savings. Although variable length frames may be used, however, the ECC block lengths are preferably fixed. Hence, whenever the data packet length is less than the ECC block length, the ECC block may be truncated (e.g., using conventional virtual zero techniques). Similar procedures may be adopted for the last block of ECC bits when the data packet is

25      larger.

As shown in the illustration, each radio data frame 40 includes a preamble 44, which is used to synchronize pseudo-random (PN) generators of the transmitter and the receiver. Link ID 46 is a field of fixed length (e.g., 16 bits long for one embodiment), and is unique to the link, thus identifying a particular subnet 10. Data from the server 12/client

5 16 is of variable length as indicated by a length field 48. Cyclic redundancy check (CRC) bits 50 may be used for error detection/correction in the conventional fashion.

For the illustrated embodiment then, each frame 52 is divided into a forward slot F, a backward slot B, a quiet slot Q and a number of radio turn around slots T. Slot F is meant for server 12-to-clients 16 communication. Slot B is time shared among a number

10 of mini-slots B<sub>1</sub>, B<sub>2</sub>, etc., which are assigned by server 12 to the individual clients 16 for their respective transmissions to the server 12. Each mini-slot B<sub>1</sub>, B<sub>2</sub>, etc. includes a time for transmitting audio, video, voice, lossy data (i.e., data that may be encoded/decoded using lossy techniques or that can tolerate the loss of some packets during transmission/reception), lossless data (i.e., data that is encoded/decoded using lossless techniques or that

15 cannot tolerate the loss of any packets during transmission/reception), low bandwidth data and/or command (Cmd.) packets. Slot Q is left quiet so that a new client may insert a request packet when the new client seeks to log-in to the subnet 10. Slots T appear between any change from transmit to receive and vice-versa, and are meant to accommodate individual radios' turn around time (i.e., the time when a half-duplex radio

20 14 switches from transmit to receive operation or vice-versa). The time duration of each of these slots and mini-slots may be dynamically altered through renegotiations between the server 12 and the clients 16 so as to achieve the best possible bandwidth utilization for the channel. Note that where full duplex radios are employed, each directional slot (i.e., F and B) may be full-time in one direction, with no radio turn around slots required.

Forward and backward bandwidth allocation depends on the data handled by the clients 16. If a client 16 is a video consumer, for example a television, then a large forward bandwidth is allocated for that client. Similarly if a client 16 is a video generator, for example a video camcorder, then a large reverse bandwidth is allocated to that 5 particular client. The server 12 maintains a dynamic table (e.g., in memory at server 12 or host 13), which includes forward and backward bandwidth requirements of all on-line clients 16. This information may be used when determining whether a new connection may be granted to a new client. For example, if a new client 16 requires more than the available bandwidth in either direction, server 12 may reject the connection request. The 10 bandwidth requirement (or allocation) information may also be used in deciding how many radio packets a particular client 16 needs to wait before starting to transmit its packets to the server 12. Additionally, whenever the channel conditions change, it is possible to increase/reduce the number of ECC bits to cope with the new channel conditions. Hence, depending on whether the information rate at the source is altered, it may require a 15 dynamic change to the forward and backward bandwidth allocation.

The use of the slotted link communication architecture described above can present challenges for accommodating asynchronous data transmissions within a subnet 10. For example, consider a file transfer process between a network master (e.g., server 12) and a client 16 that takes place in an operating environment under the control of the Windows™ 20 operating system (or one of its variants) produced by Microsoft corporation of Redmond, Washington. In such a transfer, the requesting entity (i.e., the network resource, say a personal computer, requesting the transfer of data) initiates the transfer by sending a Server Message Block (SMB) protocol command to open the file on the target platform (e.g., a server or another personal computer storing the requested material). The target devie 25 responds with another SMB command and supplies the requesting device with a pointer to

the requested file. The requesting device, using this pointer, then reads a block of N-bytes from the designated file, specifying the block size using an offset from the pointer provided by the target device. In response, the target device reads N-bytes worth of data from the subject file and delivers this information to a transmission control

- 5 protocol/internet protocol (TCP/IP) (assuming this is the transfer protocol being used) layer that handles communication between the devices. The TCP/IP layer fragments the N-bytes of data into smaller TCP/IP packets and begins transmitting the packets to the requesting device across the communication channel. As the requesting device receives the packets, it transmits acknowledgements back to the target device. After the last packet for
- 10 the N-byte transfer has been received and acknowledged, the requesting device transmits another SMB request for the next block of M-bytes from the subject file. M and N may be the same or different, and this process continues until the file transfer is complete.

The above transfer process is serialized under the control of the SMB requests and subsequent responses. This serialization is necessary because the SMB layer on the

- 15 requesting platform waits for the arrival of the last packet in the previously requested portion of the file before sending the next request. Because of this serialization, the wireless communication channel may be idle during the response times between the SMB and TCP/IP layers on both sides of the transaction. Furthermore, because the underlying wireless communication channel is a time division multiple access (TDMA) - based
- 20 architecture, any delay in these responses can cause the transmitting device to miss its allocated slot time and, hence, result in increased latency. Thus, it would be desirable to have a mechanism for avoiding such latencies in a wireless communication scheme for a computer network.

## SUMMARY OF THE INVENTION

In one embodiment, transmissions within a communication channel utilized by devices of a computer network that are outside of a device's designated time slot are accommodated through the use of a clear channel assessment time. The clear channel assessment time takes into account the device's designated transmission time slot within the communication channel with respect to those of other network devices. Thus, the clear channel assessment time may be a time period that is the product of a predetermined clear channel waiting time and a numerical representation of the difference between the device's designated transmission time slot within the communication channel and that of another network device that completed a preceding transmission. The clear channel waiting time may be specified by a network master device as part of a network connection process and the transmissions within the channel outside of a device's designated time slot may be accommodated after all regularly scheduled transmissions within the channel during a network frame period have been completed.

In another embodiment a clear channel assessment that takes into account a first device's designated transmission time slot within a communication channel with respect to those of other network devices in order to determine idle times that exist after completion of regularly scheduled transmissions within the communication channel is provided. The first device may transmit within the common communication channel upon an indication that the channel is available for transmission. Such indication is preferably made upon the expiration of a time period that is the product of a predetermined clear channel waiting time and a numerical representation of the difference between the first device's designated transmission time slot within the communication channel with respect to that of another network device. This predetermined clear channel waiting time may be designated by a network master device upon a connection thereto by the first device.

In still another embodiment, a network client having a clear channel assessment indicator and being configured to transmit within a communication channel of a computer network at a time determined in part by a notification from the clear channel assessment indicator and in part by transmission characteristics of other devices transmitting within the channel is provided. These transmission characteristics may include a numerical difference between a designated transmission slot for the network client and that of at least one of the other devices. In some cases, the channel is a time division multiplexed wireless communication channel.

In a further embodiment, a scheme for negotiating a transmission time in a time division multiplexed communication channel according to a need to transmit asynchronous data within idle times of a transmission frame period is provided. According to such a scheme, transmissions of asynchronous data within the idle times are scheduled by devices utilizing the communication channel according to a clear channel assessment time and transmission characteristics, for example designated transmission time slots within the transmission frame period, of other devices transmitting within the channel.

These and other features and advantages of the present invention will be apparent from a review of the detailed description and its accompanying drawings that follow.

## **BRIEF DESCRIPTION OF THE DRAWINGS**

The present invention is illustrated by way of example, and not limitation, in the figures of the accompanying drawings in which:

5      **Figure 1** illustrates a generalized network structure that is supported by a wireless protocol that is one embodiment of the present invention;

**Figure 2** illustrates a hierarchical arrangement for the transmission of data within a subnet according to one embodiment of the present invention; and

10     **Figure 3** illustrates a network client device configured with a collection of timers and counters, some of which may be used in accordance with one embodiment of the present invention.

## DETAILED DESCRIPTION

Described herein is a scheme for avoiding latencies in asynchronous communications with an wireless communication channel of a computer network. The present scheme is generally applicable to a variety of network environments, but finds especially useful application in a wireless computer network which is located in a home environment. Thus, the present scheme will be discussed with reference to the particular aspects of a home environment. However, this discussion should in no way be seen to limit the applicability or use of the present invention in and to other network environments and the broader spirit and scope of the present invention is recited in the claims which follow this discussion.

One important term used throughout the following discussion is "channel". As indicated above, a channel is defined as the combination of a transmission frequency (more properly a transmission frequency band) and a pseudo-random (PN) code used in a spread spectrum communication scheme. In general, a number of available frequencies and PN codes may provide a number of available channels within a subnet. Network masters and clients are capable of searching through the available channels to find a desirable channel over which to communicate with one another. Table 1 below illustrates an exemplary channel plan according to this scheme.

Table 1

Available Frequency Bands \ Available PN Codes	PN Code 1	PN Code 2	...	PN Code n
Frequency Band 1	Channel 11	Channel 12	...	Channel 1n
Frequency Band 2	Channel 21	Channel 22	...	Channel 2n
...	...	...	...	...
Frequency Band N	Channel N1	Channel N2	...	Channel Nn

In one embodiment, a channel plan using two frequency bands is adopted and details of channel selection within such a scheme is discussed in greater in the above-cited co-pending application.

As explained in greater detail in the above-cited co-pending application, when a

- 5 client device 16 join a subnet, the client receives a Connection Agreements (CAG) package from the network master device (e.g., server 12). This package includes, among other things, information regarding the forward and backward bandwidth (e.g., the slots of the channel) to which the new client 16 is entitled. In addition, the maximum number of bytes the new client 16 can send/expect in each data packet is set for each type of packet (e.g.,
- 10 video data, audio data, etc.). The Connection Agreements package may also contain information regarding the total number of data frames that the new client 16 needs to wait (i.e., before transmitting its traffic) from the start of server's transmission and the identification of the preceding client (i.e., the client that owns the preceding reverse transmission slot). The client is also assigned a unique session identifier (CS-ID).
- 15 In addition to receiving its unique CS-ID, each client is also provided (either as part of the Connections Agreement package or in another client-master exchange) with the CS-ID of the last device to transmit within a network frame 52 (i.e., the unique identifier of the last device that has an allocated time slot for transmission in the TDMA architecture of the channel). This allows each client to identify any idle times at the end of a network frame.
- 20 By knowing when these idle periods commence, clients are able to take advantage of available, but otherwise unused, channel bandwidth for asynchronous data transmissions. Such idle times may appear for a variety of reasons, for example, some devices in the subnet 10 may not always make full use of their available transmission times. Also, the entire network frame period may be more than is needed to accommodate existing devices
- 25 within the subnet 10, thus providing excess time at the end of such a period.

After receiving the Connection Agreements packet, the client 16 configures itself to transmit its data in its assigned time slot (e.g., B<sub>1</sub>, B<sub>2</sub>, etc.) and waits for that slot to come around. At the designated time slot, the client 16 may initiate normal communications with the server 12 and transmit any data or commands it may have. In order to help

5 maintain proper time slot synchronization, part of this configuration process may involve programming an accurate slot timer (AST).

**Figure 3** illustrates a client 16 configured with an AST 54. The AST 54 may be a conventional timer (e.g., a register or memory location, which is incremented or decremented on a regular basis according to an appropriate clock signal maintained by the

10 client device) that is programmed using information provided by the network master as part of the connection process. For example, the network master (e.g., server 12) may provide the client 16 with an indication of where it lies in within the slotted link structure of the communication protocol (i.e., at what time it should begin its transmissions). By then monitoring this time using AST 54 (which may be reset upon each master

15 transmission or each transmission of the associated client), the client 16 can accurately predict when it should begin its transmission. Of course, the expiration of an AST-monitored time should not allow transmission in the event a client detects that a preceding client has not yet completed its own transmissions. Otherwise, the transmissions may overlap and cause confusion within the network.

20 In conjunction with the use of an AST 54, a client 16 may also incorporate an early trigger timer (ETT) 55 to ensure that packets are ready for transmission within the subnet when the client's transmission slot arrives. That is, an ETT 55 may be used to advance the internal construction of a packet or packets for transmission, with the goal being to have those packets assembled and ready for transmission when the client's transmission slot 25 becomes available. Thus the ETT 55 (which may be implemented as a conventional timer)

triggers the process of formation of packets and their error protection bits, etc., using a packet construction engine 56 (which in some cases may be a part of client 16 or in other cases may be a part of radio 14) to keep a few packets, at least, prepared before the actual start of transmission. For example, in one embodiment packets could be assembled and

5 stored for transmission one network frame in advance of their actual transmission. Such preassembly is helpful in avoiding the idle times on the channel when the first few packets for a transmission sequence are being formed. Although collecting data one network frame in advance of transmission may penalize the system with a one network frame latency at the first transmission slot, it is expected that this period can be reduced to a few

10 milliseconds.

As indicated above, each client 16 may be required to keep track of the present client occupying the channel, thereby trying to detect its immediately preceding client in line. If the channel is quiet, the current client waits for a predetermined length of time before starting its own transmission. In one embodiment, the length of this waiting time

15 depends on the quiet time threshold allowed between two clients (termed the clear channel assessment or CCA time) and the number of clients yet to transmit before the current client. For example, the waiting time may be the product of the CCA threshold and the number of devices yet to transmit. This waiting time calculation thus makes use of the order of transmission that is established during the connection setup. The only exception

20 to the quiet time is the Q slot, when all on-line clients 16 should refrain from transmitting.

As indicated, a CCA transmission (i.e., a client's transmission based on a waiting period timeout) is essentially the result of a decision made by the client that the channel is free and hence is suitable for transmission. Each client device keeps track of how many devices have completed the transmission from the time the network master completed its

25 transmission. This can be done using a counter that is incremented each time a client

transmission is detected and the implementation of such counters is well known in the art. Then, by knowing the index of its own position in the network frame (a value received from the network master at the time a connection is established), each client device can determine when to initiate a CCA transmission. This is best illustrated with an example.

5        Each client device may program an associated CCA timer 57 (see Figure 4) to a predetermined value. The network master may specify this value at the time a master-client connection is established and it generally represents a period of time that must expire before a network client is permitted to assume that another client is not using the channel.

Now, suppose a client device is 5<sup>th</sup> in line for transmission after the master device and

10      suppose it detects a clear channel (e.g., because its CCA timer 57 times out) after the device that is second in line has completed its transmission. Because the device is 5<sup>th</sup> in line, it cannot immediately begin its own transmissions (after all, there are still two client devices in line ahead of it). However, the 5<sup>th</sup> device may increment an associated CCA counter 58 at this time. If then both the 3<sup>rd</sup> and 4<sup>th</sup> devices are silent (i.e., if the 5<sup>th</sup> device's

15      CCA timer 57 times out two more times in succession), then the 5<sup>th</sup> device will have again twice incremented its CCA counter 58 and may immediately start its own transmission on the 3<sup>rd</sup> CCA detection.

Note that in the above example, the 3<sup>rd</sup> device in line will also have detected the absence of the second device's transmission and may therefore immediately start its own

20      transmission if indeed it has traffic to send. In other words, the 3<sup>rd</sup> device's waiting period is only one CCA timeout from the absence of the 2<sup>nd</sup> device's transmissions, while that of the 5<sup>th</sup> device is three CCA timeout periods. Similarly, the 4<sup>th</sup> device in line is two CCA timeout periods away from the second device. This use of the client slot assignments in determining when CCA transmissions may be initiated allows the client devices to

maintain their relative sequence with one another within the slotted link structure of the communication channel.

The CCA timeout period may also be used to detect idle times at the end of a network frame 52. For example, by knowing the CS-ID of the last device to transmit within the TDMA architecture, all clients 16 can determine when the idle time commences by monitoring the transmission of the client that is last in line. All clients 16 in the subnet 10 can then estimate the duration of the idle time after the last-in-line client has completed its transmission (by knowing the total available time for a network frame, which as indicated above, has a fixed duration). By sharing this available idle time amongst themselves, the clients 16 of the subnet 10 can provide for asynchronous data transmissions in the subnet 10. Of course, other methods of detecting idle times may be used.

The idle time "sharing" plan between the clients makes use of the CCA in that each client waits a time  $T_{idle} = T_{CCA} * C$  (microseconds) before transmitting a packet in the idle time.  $T_{CCA}$  may be the standard CCA time for a regular transmission and  $C$  may be

determined as the difference between the transmission slot number of the current client and that of the client from which the immediately preceding packet was received (or monitored). If the master device has any asynchronous data to transmit, it may use  $C = 1$  and begin its transmission from the time of the end of the transmission of the last-in-line client.

In order to provide a somewhat fair allocation scheme, in one embodiment each client is permitted to transmit only one packet in the idle time at the end of a network frame 52. This allows the devices in subnet 10 to take turns transmitting asynchronous (e.g., low priority) data over the channel. This protocol may be abandoned and a transmission commenced if a packet from a previous client is detected in the idle time and there is sufficient time in the idle period for a packet transmission. However, before any

transmission in the idle time, a device wishing to send data should allow sufficient time for the Q slot before the commencement of the next network frame.

Within the present scheme, there are some idle-time transmission situations that warrant further discussion. For example, consider the situation of a short idle time.

- 5 Because the time at which a device may transmit in the idle time is determined by transmission slot number (C), it is possible that only a subset of the total number of network devices get the opportunity to make use of the idle time in each network frame period. To make idle time allocation more fair, the present scheme allows devices in a subnet 10 to request reallocation of its transmission slot by sending such a request to the
- 10 network master. Thus, if a client determines that it has a significant amount of asynchronous traffic to send but that it is unable to do so because of its transmission slot number, that client can request a new transmission slot, earlier in a network frame period, so as to have a better chance of making use of the idle time at the end of the frame.

If a device wishes to transmit multiple packets within the idle time of a single network frame, a two-step process may be invoked. First, the device transmits an initial packet in its regular space in the idle time, determined according to the above protocol. Then, the device may reprogram its  $T_{idle}$  time such that  $T_{idle} = T_{CCA} * N$ , where N is the total number of devices (including the master) in the subnet 10. The next idle time transmission for the subject device can then occur at this new  $T_{idle}$  time, provided sufficient time remains before the Q slot. If a packet from another device is received before this new  $T_{idle}$  time expires, the CCA may be reprogrammed to the time between the device's transmission and the reception of the newly received packet. This helps assure an equal opportunity for all devices in the subnet 10 to use the idle time for transmission of asynchronous data.

When a subnet 10 is sharing a channel with one or more overlapping subnets, the

- 25 start and end of each subnet's transmission periods should be strictly limited to a time

duration within each network frame 52. If there are idle times within such a duration, the devices in the subnets may make use of the time in the manner described above, however, such transmissions should not exceed the time duration of their respective subnet's transmission window.

5        During channel change situations (described in detail in the above-cited co-pending application), client devices are prohibited from transmitting until permitted to do so by the master device or certain time outs. This protocol should be strictly observed and idle time transmissions halted during such periods so that there is no loss of data within the subnet 10.

10      Packet size will play a role in determining how many devices are afforded the opportunity to transmit within an idle time. That is, even if a device's time for transmission in an idle time has arrived, that device should not transmit if it detects an ongoing transmission of another device. Also, devices sending smaller packets during an idle time may be benefited to a lesser extent than those transmitting larger packets, as a 15 device is only (usually) able to transmit one packet per idle time. To balance this equation, and reduce the potential for collisions within an idle time, one embodiment of the present scheme restricts devices to a predetermined time duration for any such transmissions. In other words, only packets up to a predetermined size are eligible for transmission in an idle time.

20      Thus, a scheme for synchronizing communications within a computer network communication channel has been described. Although discussed with reference to certain illustrated embodiments, the present invention should not be limited thereby. Instead, the present invention should only be measured in terms of the claims that follow.

## CLAIMS

What is claimed is:

- 1 1. A method, comprising accommodating within a common communication channel having
- 2 designated transmission time slots for various devices of a computer network transmissions
- 3 within the channel outside of a device's designated time slot through the use of a clear
- 4 channel assessment time.
  
- 1 2. The method of claim 1 wherein the clear channel assessment time takes into account the
- 2 device's designated transmission time slot within the communication channel with respect
- 3 to those of other network devices.
  
- 1 3. The method of claim 2 wherein the clear channel assessment time comprises a time
- 2 period that is the product of a predetermined clear channel waiting time and a numerical
- 3 representation of the difference between the device's designated transmission time slot
- 4 within the communication channel and that of another network device that completed a
- 5 preceding transmission.
  
- 1 4. The method of claim 3 wherein the clear channel waiting time is specified by a network
- 2 master device as part of a network connection process.
  
- 1 5. The method of claim 1 wherein the transmissions within the channel outside of a
- 2 device's designated time slot are accommodated after all regularly scheduled transmissions
- 3 within the channel during a network frame period have been completed.

1 6. A method, comprising maintaining a clear channel assessment that takes into account a  
2 first device's designated transmission time slot within a communication channel with  
3 respect to those of other network devices in order to determine idle times that exist after  
4 completion of regularly scheduled transmissions within the communication channel.

1 7. The method of claim 6 wherein the first device transmits within the common  
2 communication channel upon an indication that the channel is available for transmission.

1 8. The method of claim 7 wherein the indication is made upon the expiration of a time  
2 period that is the product of a predetermined clear channel waiting time and a numerical  
3 representation of the difference between the first device's designated transmission time slot  
4 within the communication channel with respect to that of another network device.

1 9. The method of claim 8 wherein the predetermined clear channel waiting time is  
2 designated by a network master device upon a connection thereto by the first device.

1 10. A network client comprising a clear channel assessment indicator and configured to  
2 transmit within a communication channel of a computer network at a time determined in  
3 part by a notification from the clear channel assessment indicator and in part by  
4 transmission characteristics of other devices transmitting within the channel.

1 11. The network client of claim 10 wherein the transmission characteristics comprise a  
2 numerical difference between a designated transmission slot for the network client and that  
3 of at least one of the other devices.

1 12. The network client of claim 10 wherein the channel is a time division multiplexed  
2 wireless communication channel.

1 13. A method comprising negotiating a transmission time in a time division multiplexed  
2 communication channel independent of a need to transmit asynchronous data within idle  
3 times of a transmission frame period.

1 14. The method of claim 13 wherein transmissions of asynchronous data within the idle  
2 times are scheduled by devices utilizing the communication channel according to a clear  
3 channel assessment time and transmission characteristics of other devices transmitting  
4 within the channel.

1 15. The method of claim 14 wherein the transmission characteristics comprise designated  
2 transmission time slots within the transmission frame period.

1 16. A method comprising accommodating asynchronous data transmissions within a  
2 synchronized network in which inter-node communications are organized into frames of  
3 time periods by permitting such asynchronous communications within otherwise idle times  
4 within the frames.

1 17. The method of claim 16 wherein use of the otherwise idle times within the frames takes  
2 into account a transmitting node's designated transmission time within a particular frame  
3 with respect to transmission times of other nodes of the network.

1 18. The method of claim 16 wherein the asynchronous data transmissions are self-  
2 organized and/or self-synchronized by nodes of the network without direct scheduling  
3 assistance from a network master.

## ABSTRACT

Transmissions within a communication channel utilized by devices of a computer network that are outside of a device's designated time slot are accommodated through the use of a clear channel assessment time. The clear channel assessment time takes into 5 account the device's designated transmission time slot within the communication channel with respect to those of other network devices. Thus, the clear channel assessment time may be a time period that is the product of a predetermined clear channel waiting time and a numerical representation of the difference between the device's designated transmission time slot within the communication channel and that of another network device that 10 completed a preceding transmission. The clear channel waiting time may be specified by a network master device as part of a network connection process and the transmissions within the channel outside of a device's designated time slot may be accommodated after all regularly scheduled transmissions within the channel during a network frame period have been completed.

22-141 50 SHEETS  
22-142 100 SHEETS  
22-144 200 SHEETS

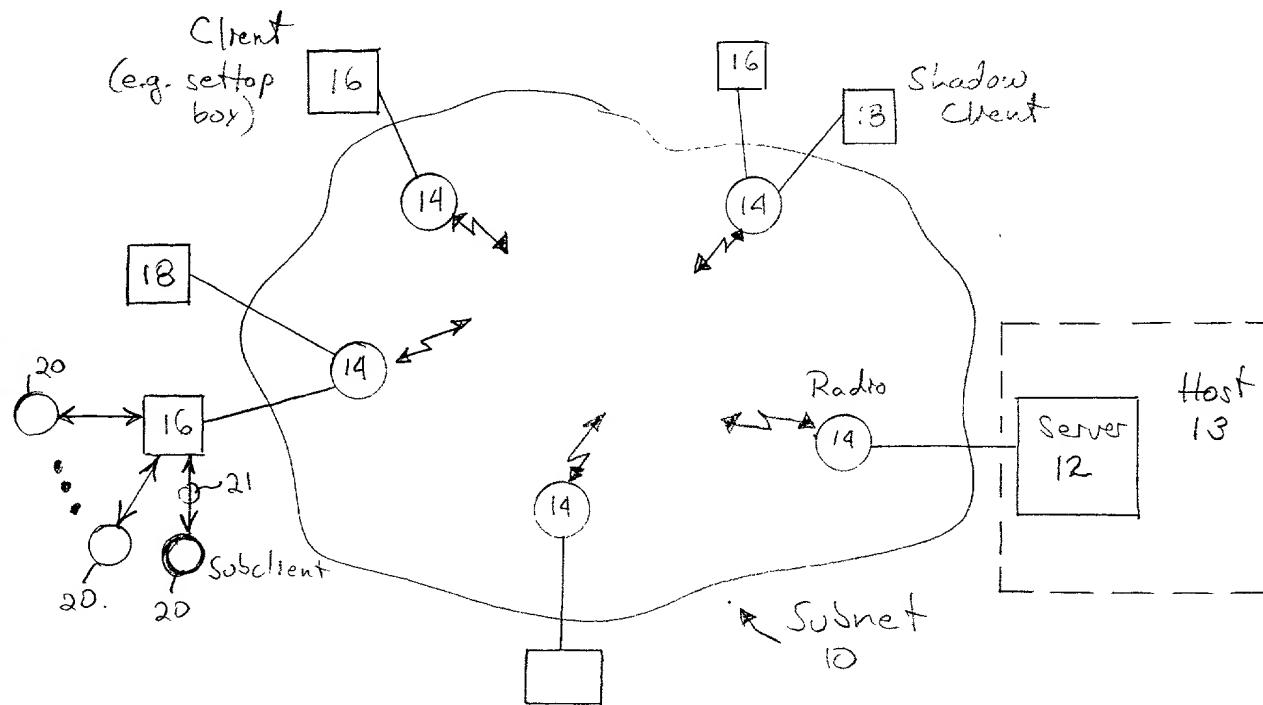


Fig. 1

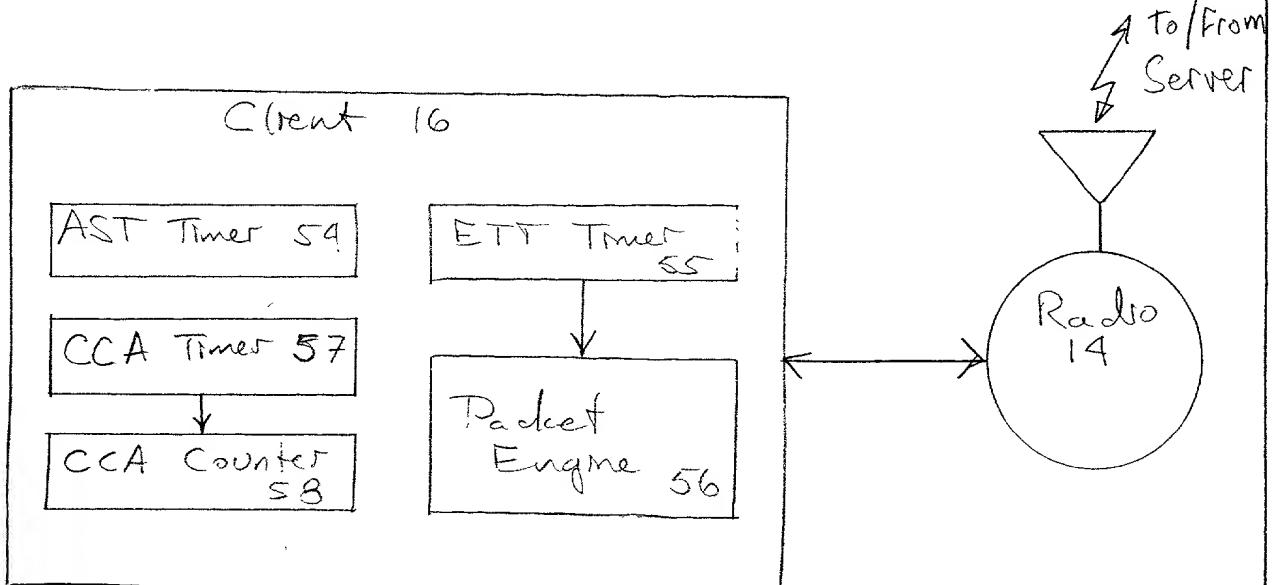


Fig. 3

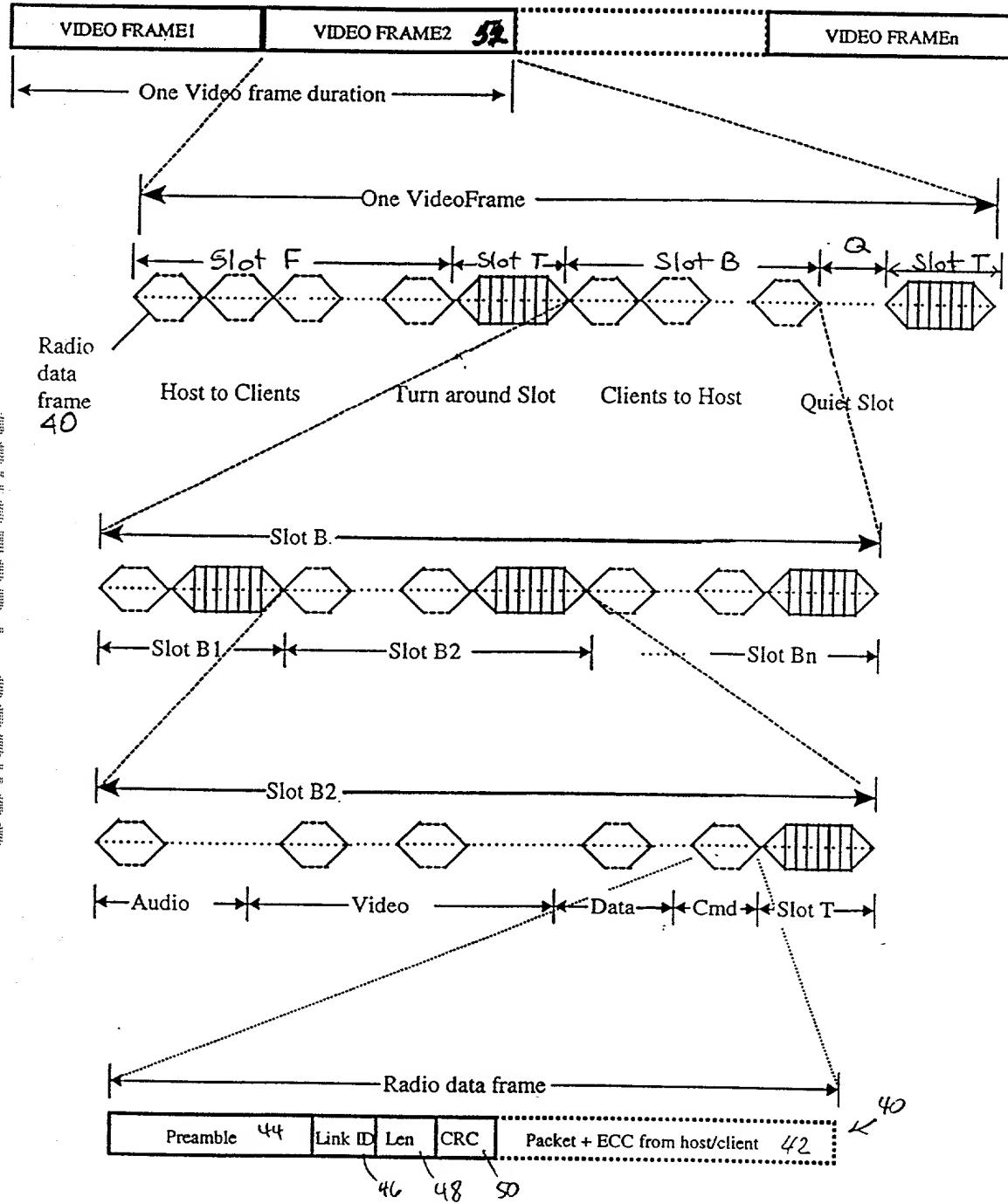


Figure 2

Attorney's Docket No.: 003498.P033

Patent

**DECLARATION AND POWER OF ATTORNEY FOR PATENT APPLICATION**

As a below named inventor, I hereby declare that:

My residence, post office address and citizenship are as stated below, next to my name.

I believe I am the original, first, and sole inventor (if only one name is listed below) or an original, first, and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled

**METHOD AND APPARATUS FOR ACCOMMODATING ASYNCHRONOUS DATA TRANSMISSIONS IN A WIRELESS COMPUTER NETWORK**

the specification of which

X is attached hereto.  
— was filed on \_\_\_\_\_ as  
United States Application Number \_\_\_\_\_  
or PCT International Application Number \_\_\_\_\_  
and was amended on \_\_\_\_\_  
(if applicable)

I hereby state that I have reviewed and understand the contents of the above-identified specification, including the claim(s), as amended by any amendment referred to above. I do not know and do not believe that the claimed invention was ever known or used in the United States of America before my invention thereof, or patented or described in any printed publication in any country before my invention thereof or more than one year prior to this application, that the same was not in public use or on sale in the United States of America more than one year prior to this application, and that the invention has not been patented or made the subject of an inventor's certificate issued before the date of this application in any country foreign to the United States of America on an application filed by me or my legal representatives or assigns more than twelve months (for a utility patent application) or six months (for a design patent application) prior to this application.

I acknowledge the duty to disclose all information known to me to be material to patentability as defined in Title 37, Code of Federal Regulations, Section 1.56.

I hereby claim foreign priority benefits under Title 35, United States Code, Section 119(a)-(d), of any foreign application(s) for patent or inventor's certificate listed below and have also identified below any foreign application for patent or inventor's certificate having a filing date before that of the application on which priority is claimed:

<u>Prior Foreign Application(s)</u>	<u>Priority Claimed</u>
(Number) _____	(Country) _____
(Day/Month/Year Filed)	Yes _____ No _____
(Number) _____	(Country) _____
(Day/Month/Year Filed)	Yes _____ No _____
_____	_____

(Number)

(Country)

(Day/Month/Year Filed)

Yes No

I hereby claim the benefit under title 35, United States Code, Section 119(e) of any United States provisional application(s) listed below

(Application Number)

Filing Date

(Application Number)

Filing Date

I hereby claim the benefit under Title 35, United States Code, Section 120 of any United States application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code, Section 112, I acknowledge the duty to disclose all information known to me to be material to patentability as defined in Title 37, Code of Federal Regulations, Section 1.56 which became available between the filing date of the prior application and the national or PCT international filing date of this application:

09/151,579

September 11, 1998

Pending

(Application Number)

Filing Date

(Status -- patented,  
pending, abandoned)

(Application Number)

Filing Date

(Status -- patented,  
pending, abandoned)

I hereby appoint Farzad E. Amini, Reg. No. P42,261; Aloysius T. C. AuYeung, Reg. No. 35,432; Amy M. Armstrong, Reg. No. 42,265; William Thomas Babbitt, Reg. No. 39,591; Carol F. Barry, Reg. No. 41,600; Jordan Michael Becker, Reg. No. 39,602; Bradley J. Bereznak, Reg. No. 33,474; Michael A. Bernadicou, Reg. No. 35,934; Roger W. Blakely, Jr., Reg. No. 25,831; Gregory D. Caldwell, Reg. No. 39,926; Yong S. Choi, Reg. No. P43,324; Thomas M. Coester, Reg. No. 39,637; Michael Anthony DeSanctis, Reg. No. 39,957; Daniel M. De Vos, Reg. No. 37,813; Robert Andrew Diehl, Reg. No. 40,992; Tarek N. Fahmi, Reg. No. 41,402; James Y. Go, Reg. No. 40,621; Dinu Gruia, Reg. No. 42,996; Thomas A. Hassing, Reg. No. 36,159; Phuong-Quan Hoang, Reg. No. 41,839; Willmore F. Holbrow III, Reg. No. 41,845; George W Hoover II, Reg. No. 32,992; Eric S. Hyman, Reg. No. 30,139; Dag H. Johansen, Reg. No. 36,172; William W. Kidd, Reg. No. 31,772; Michael J. Mallie, Reg. No. 36,591; Andre L. Marais, under 37 C.F.R. § 10.9(b); Paul A. Mendonsa, Reg. No. 42,879; Darren J. Milliken, Reg. 42,004; Thinh V. Nguyen, Reg. No. 42,034; Kimberley G. Nobles, Reg. No. 38,255; Daniel E. Ovanezian, Reg. No. 41,236; Babak Redjaian, Reg. No. 42,096; James H. Salter, Reg. No. 35,668; William W. Schaal, Reg. No. 39,018; James C. Scheller, Reg. No. 31,195; Anand Sethuraman, Reg. No. 43,351; Charles E. Shermwell, Reg. No. 40,171; Maria McCormack Sobrino, Reg. No. 31,639; Stanley W. Sokoloff, Reg. No. 25,128; Judith A. Szepesi, Reg. No. 39,393; Vincent P. Tassinari, Reg. No. 42,179; Edwin H. Taylor, Reg. No. 25,129; Glenn E. Von Tersch, Reg. No. 41,364; George G. C. Tseng, Reg. No. 41,355; Lester J. Vincent, Reg. No. 31,460; John Patrick Ward, Reg. No. 40,216; Stephen Warhola, Reg. No. 43,237; Charles T. J. Weigell, Reg. No. 43,398; Ben J. Yorks, Reg. No. 33,609; and Norman Zafman, Reg. No. 26,250; my attorneys, of BLAKELY, SOKOLOFF, TAYLOR & ZAFMAN LLP, with offices located at 12400 Wilshire Boulevard, 7th Floor, Los Angeles, California 90025, telephone (408) 720-8300, and James R. Thein, Reg. No. 31,710, my patent attorney; with full power of substitution and revocation, to prosecute this application and to transact all business in the Patent and Trademark Office connected herewith.

Send all correspondence to Tarek N. Fahmi, Reg. No. 41,402, BLAKELY, SOKOLOFF, TAYLOR & ZAFMAN LLP, 12400 Wilshire Boulevard, 7th Floor, Los Angeles, California 90025, telephone (408) 720-8300.

I declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of this application or any patent issuing thereon.

Full Name of Sole/First Inventor Rajugopal R. Gubbi

Inventor's Signature Rajugopal R. Gubbi Date 23<sup>rd</sup> Sept 1999

Residence Fair Oaks, California (City, State) Citizenship India (Country)

Post Office Address 8842 Winding Way, Apt # 123  
Fair Oaks, California 95628-6467

Title 37, Code of Federal Regulations, Section 1.56  
Duty to Disclose Information Material to Patentability

(a) A patent by its very nature is affected with a public interest. The public interest is best served, and the most effective patent examination occurs when, at the time an application is being examined, the Office is aware of and evaluates the teachings of all information material to patentability. Each individual associated with the filing and prosecution of a patent application has a duty of candor and good faith in dealing with the Office, which includes a duty to disclose to the Office all information known to that individual to be material to patentability as defined in this section. The duty to disclosure information exists with respect to each pending claim until the claim is cancelled or withdrawn from consideration, or the application becomes abandoned. Information material to the patentability of a claim that is cancelled or withdrawn from consideration need not be submitted if the information is not material to the patentability of any claim remaining under consideration in the application. There is no duty to submit information which is not material to the patentability of any existing claim. The duty to disclosure all information known to be material to patentability is deemed to be satisfied if all information known to be material to patentability of any claim issued in a patent was cited by the Office or submitted to the Office in the manner prescribed by §§1.97(b)-(d) and 1.98. However, no patent will be granted on an application in connection with which fraud on the Office was practiced or attempted or the duty of disclosure was violated through bad faith or intentional misconduct. The Office encourages applicants to carefully examine:

(1) Prior art cited in search reports of a foreign patent office in a counterpart application, and

(2) The closest information over which individuals associated with the filing or prosecution of a patent application believe any pending claim patentably defines, to make sure that any material information contained therein is disclosed to the Office.

(b) Under this section, information is material to patentability when it is not cumulative to information already of record or being made or record in the application, and

(1) It establishes, by itself or in combination with other information, a prima facie case of unpatentability of a claim; or

(2) It refutes, or is inconsistent with, a position the applicant takes in:

(i) Opposing an argument of unpatentability relied on by the Office, or

(ii) Asserting an argument of patentability.

A prima facie case of unpatentability is established when the information compels a conclusion that a claim is unpatentable under the preponderance of evidence, burden-of-proof standard, giving each term in the claim its broadest reasonable construction consistent with the specification, and before any consideration is given to evidence which may be submitted in an attempt to establish a contrary conclusion of patentability.

(c) Individuals associated with the filing or prosecution of a patent application within the meaning of this section are:

(1) Each inventor named in the application;

(2) Each attorney or agent who prepares or prosecutes the application; and

(3) Every other person who is substantively involved in the preparation or prosecution of the application and who is associated with the inventor, with the assignee or with anyone to whom there is an obligation to assign the application.

(d) Individuals other than the attorney, agent or inventor may comply with this section by disclosing information to the attorney, agent, or inventor.